Contingent Allegiance (CA) condition is defined as follows in section 6.6 of the Standard:

"The contingent allegiance condition shall exist following the return of a CHECK CONDITION or COMMAND TERMINATED status."

For cases involving a queue with commands from multiple initiators and a CA existing for one initiator, there are conflicting statements in the Standard in sections 6.6 and 6.8.2. Specifically, section 6.6 states:

"Execution of all tagged I/O processes for the I_T_L nexus for which the contingent allegiance exists shall be suspended until the contingent allegiance condition is cleared."

This implies that only the commands for the I_T_L nexus which is in a CA be suspended. Section 6.8.2 states:

"The first method allows the execution of tagged I/O processes to resume when the contingent allegiance or extended contingent allegiance condition is cleared. The target returns BUSY status to other initiators while the contingent allegiance or extended contingent allegiance condition exists. During this time, all tagged I/O processes are suspended."

This appears to be a direct contradiction of the earlier statement.