

Introduction to Thin Provisioning Proposal

(and FLASH pre-write)

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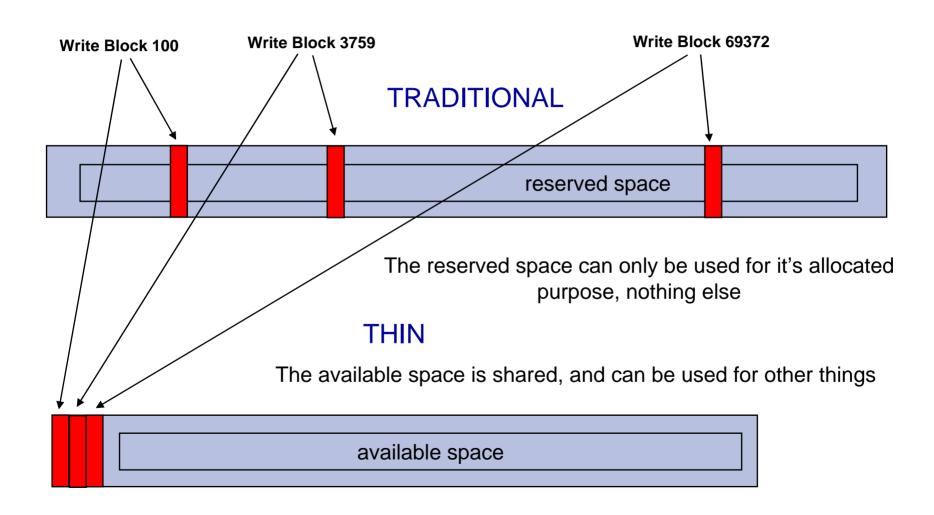
Thin Provisioning – What is it?

- Traditional storage devices reserve physical storage for every possible disk address (block) whether it contains useful data or not.
- ▶ Hosts rarely use 100% of possible addresses

- "Thin" devices allocate physical storage only for those blocks that contain useful data.
 - Useful data is data that has been written.
- The Storage is virtualized



Storing LUN data





Thin Provisioning – Why?

- "Thin" devices require less real H/W
- "Thin" devices require less space
- "Thin" devices require less power
- "Thin" devices substantially reduce TCO

 Storage virtualization brings substantial benefits to customers just like host virtualization



Thin Provisioning – Opportunity

Deleting a file – reusing deallocated space



Since the space is already reserved deleting a file has no impact on utilization

THIN

Reclaiming the now unused space is now a desired action



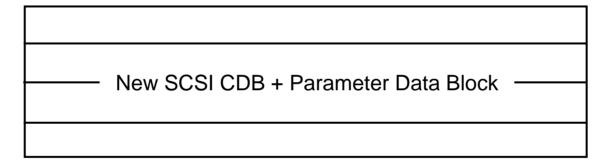


Thin Provisioning – Issues

- Today, hosts assume the traditional model
- Hosts don't tell storage when the data in a block is useful and when it is not (such as when a file is deleted from a file system).
- Storage has no standard way to inform the host of the provisioning status changes.
- Need to create standard APIs for Thin Provisioning

Standard APIs

- Create a new SCSI Command for Hole Punch/Delete/Pre-Write Function
 - Block numbers not in CDB (in parameter data block)



LBA
Length
LBA
Length
Length

Settable Thresholds

- Soft Threshold like tape programmable early warning Begin reporting to host (recovered error)
- Hard Threshold like tape EOM Thin space ran out Writes start to fail (hard error)
- Mode Page readable/settable?
- Query state of Thin/Thick status of block:
 - Send LBA/allocation length returns bit/byte map



Report Supported Operation Codes

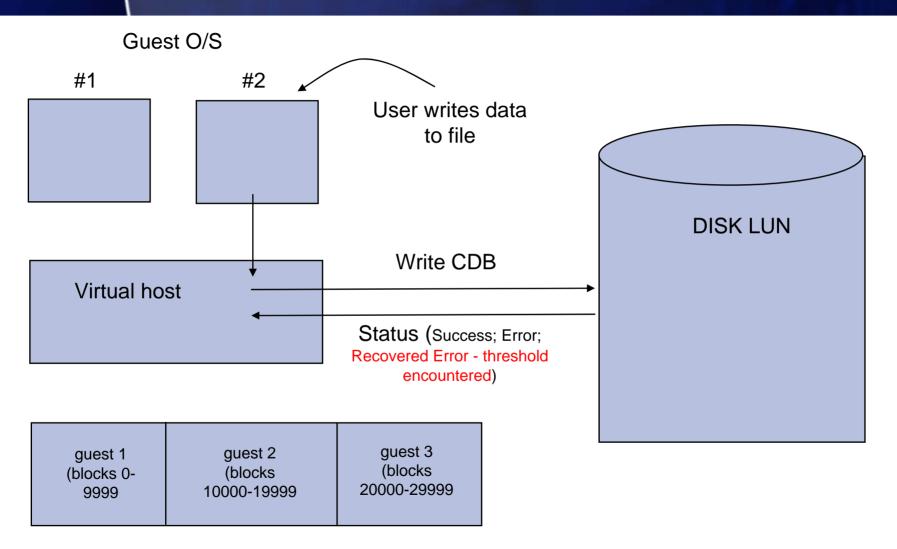
Tells host new command exists

► Create new SCSI return codes (samples)

- CC: RECOVERED ERROR + SOFT CAPACITY THRESHOLD ENCOUNTERED
- CC: RECOVERED ERROR + HARD CAPACITY THRESHOLD ENCOUNTERED
- CC: HARDWARE ERROR/DATA PROTECT + HARD CAPACITY THRESHOLD ENCOUNTERED

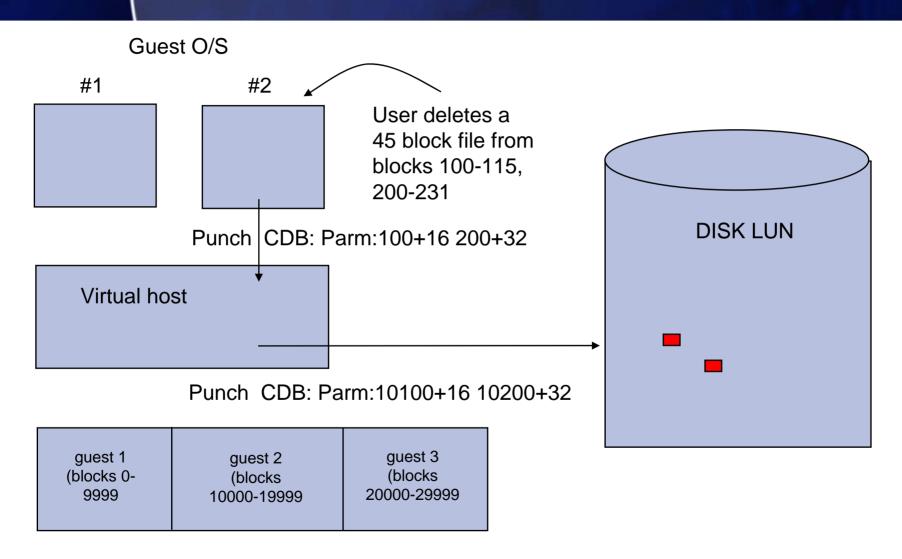


Use Cases



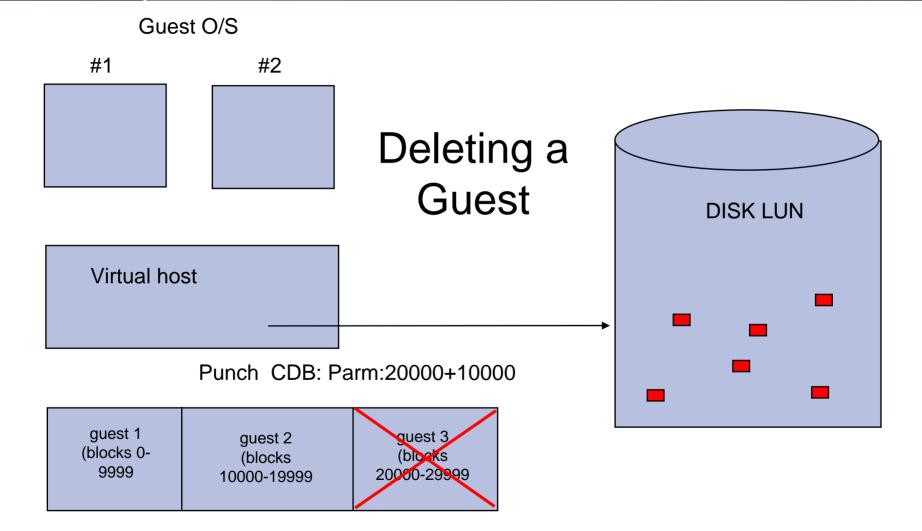
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Use Cases



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Use Cases



- ▶ Flash Memory "disks"
 - Read = 20us
 - Write (to erased area) = 200us
 - Write (to non-erased area) = 1.7ms
 - Erase = 1.5ms
- Erase on delete will speed up next write.
 - Perfect use for Punch/Delete/Pre-Write CDB



DISCUSSION