

- Training on each I/O **limits performance**
 - Keep option open to train less often

- Eliminating QAS **limits performance**
 - Keep QAS for those target / initiator pairs that want to go fast

- Incremental changes to SPI-4 eventually add up

Performance impact depends upon :

- I/O size → • Significant for ≤ 8 KB I/Os
- I/O direction (read / write) → • Higher impact for Write on larger transfers
- Queue Depth → • Impact increases as queue depth < 8
- Target head movement → • Impact increases as head movement decreases

Do we create a specification that optimizes a specific operational environment, or one that can work well over a wide range of environments ?

- Evaluate:
 - P + QAS baseline Packetized, QAS, no training
 - P + QAS + Training Packetized, QAS, 2us training
 - P + Training Packetized, no QAS, 2 us training

- This analysis does not consider Parallel option, with or without training

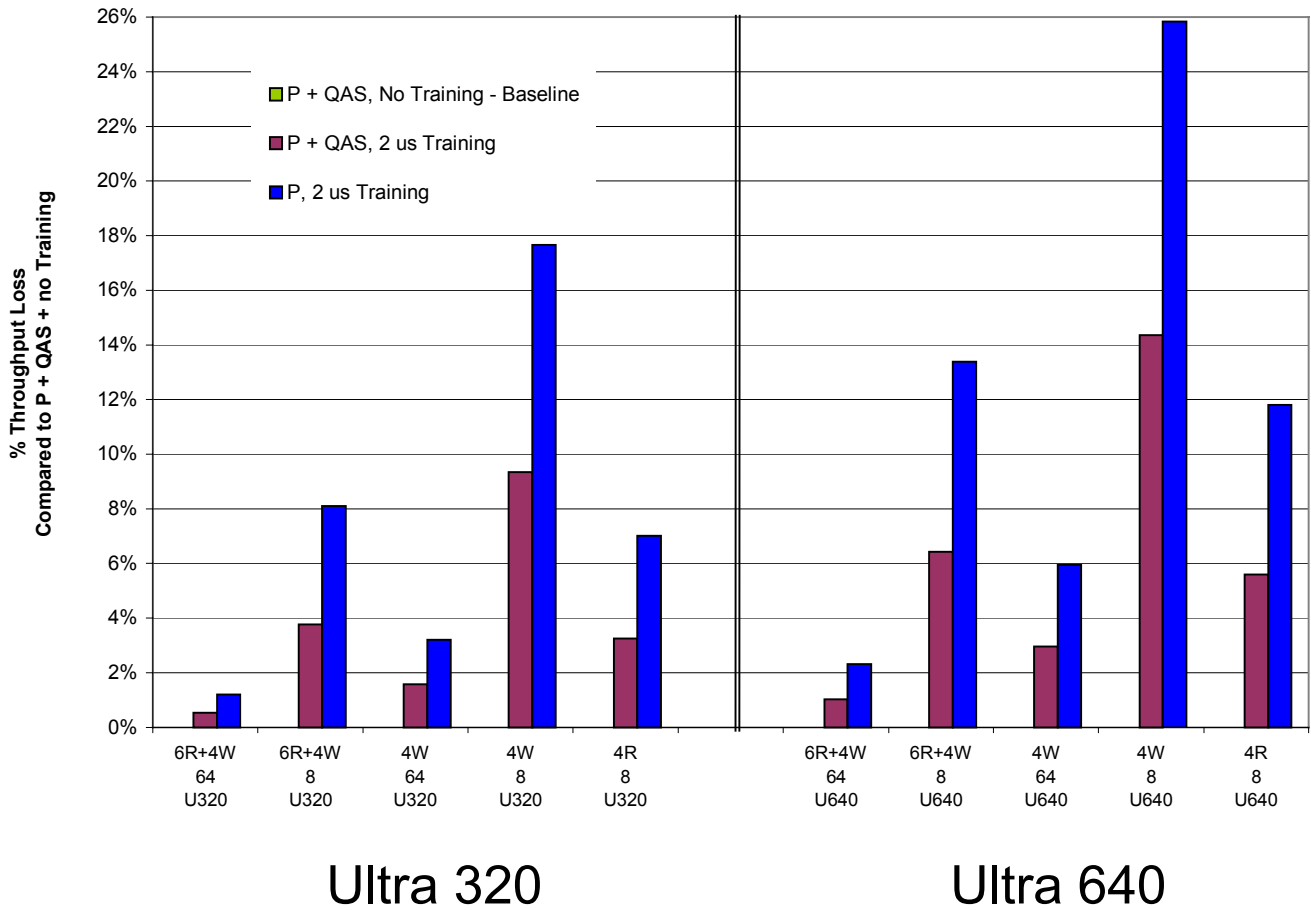
Training & QAS impact < 2% at 320 at large blocks, at 10 Q'd I/Os

Impact of Training 9% at 8KB, 4 Q'd

Impact of QAS another 9% at 8KB, 4 Q'd

Impact of each Training & QAS is 14% each at Ultra640 - it gets worse!

Performance Cost of Changes



I/O size in Kbytes

4W = 4 Queued Write I/Os

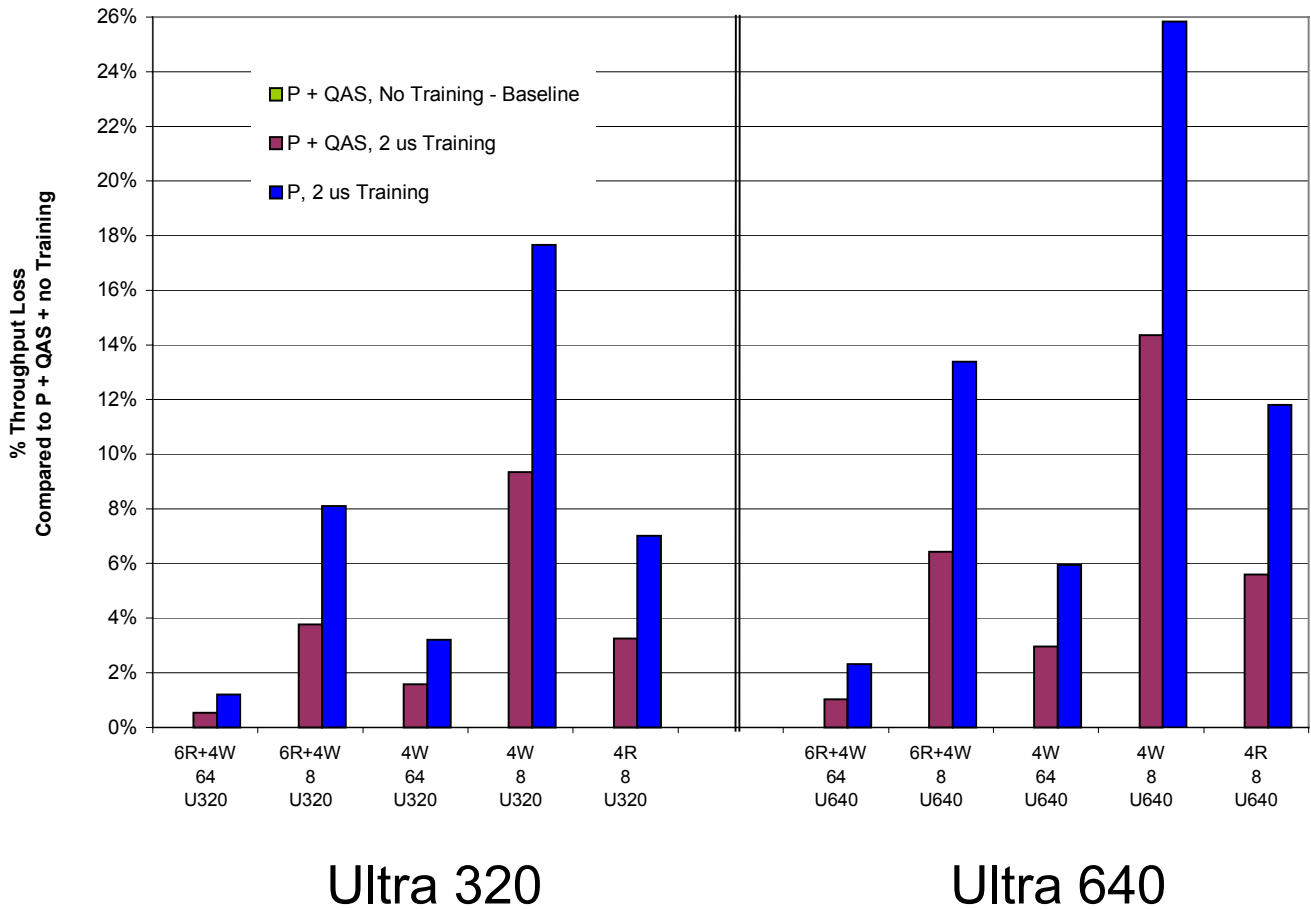
Impact of changes is small at Q Depth > 8. Today Q Depth = 1 for sequential. Which applications will change the way they issue I/Os for Ultra320?

Impact of Training & QAS each are significant at smaller Queue depths, together they can cripple SCSI

Lots of I/O happens with blocks of 4K to 8KB

These problems get worse at Ultra640

Performance Cost of Changes



I/O size in Kbytes

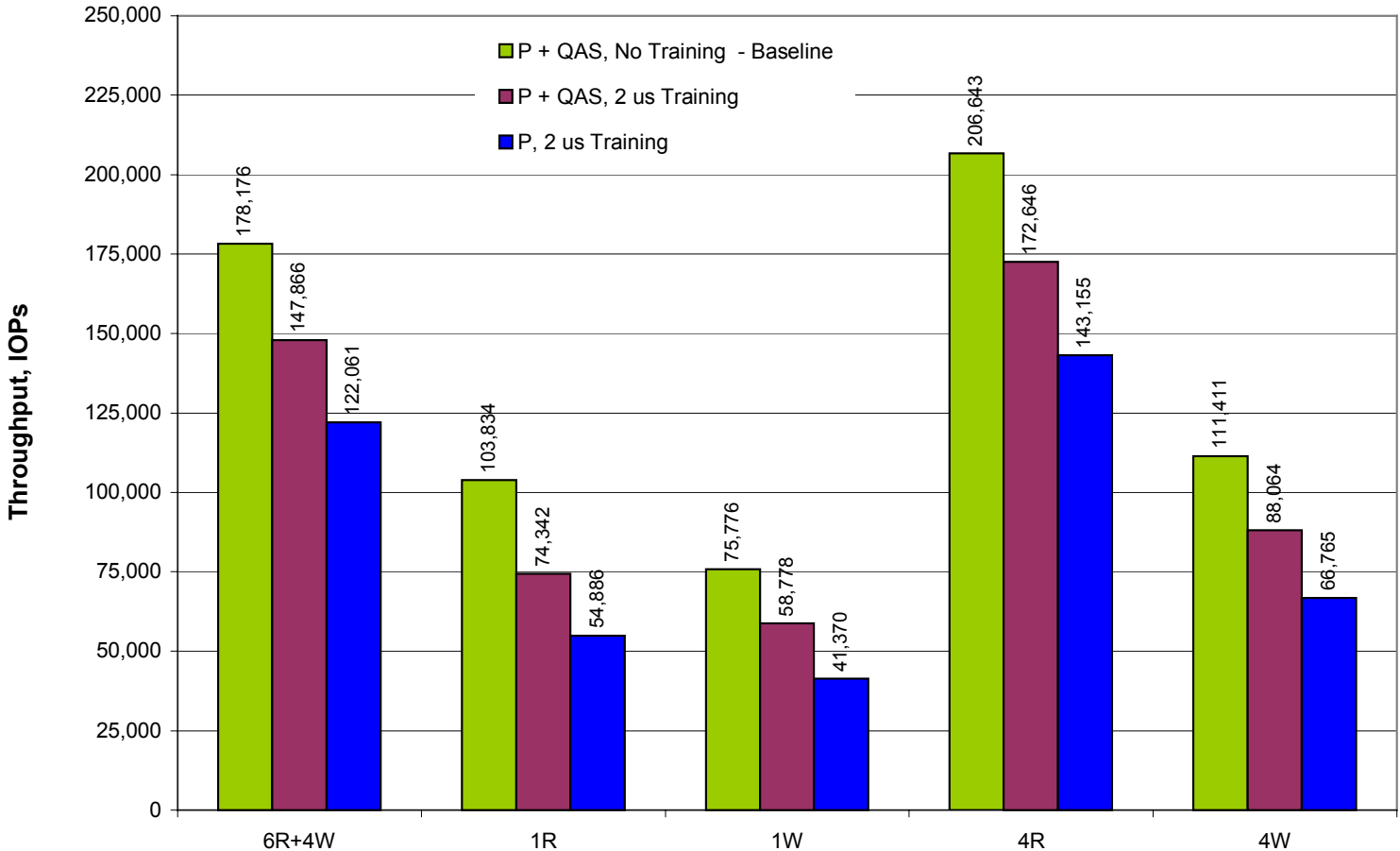
4W = 4 Queued Write I/Os

Theoretical results - only counts SCSI protocol timings

Realistic IOPs values are 40 - 50 % lower than these theoretical results at 0.5 KB

Ultra320 with Packetized & training will be slower than Ultra160 Non Packetized at 1 Q'd I/O per drive - same test as we use today !

Ultra320 Performance Cost of Changes, IOPs, 0.5 KB I/Os

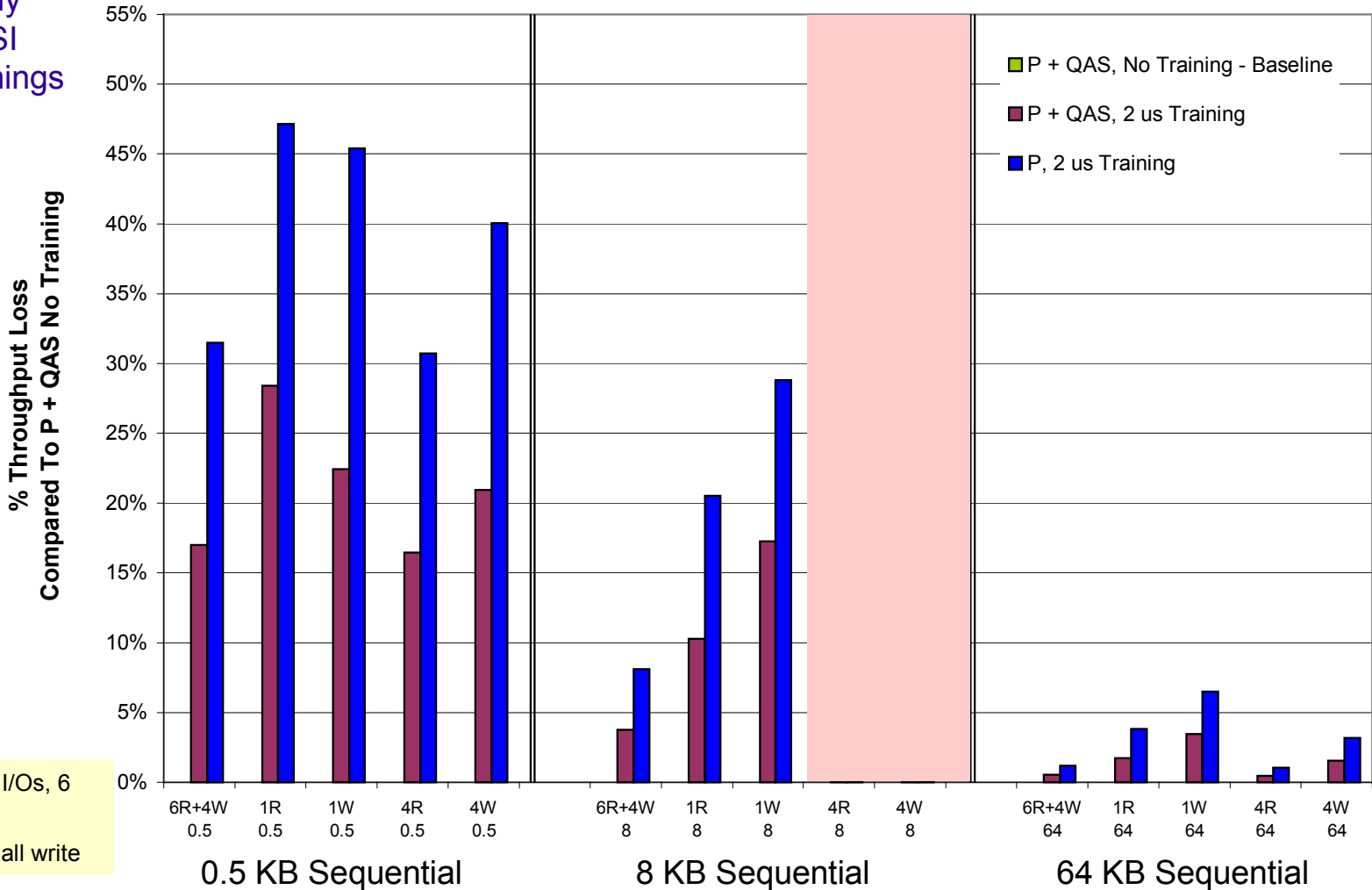


6R+4W = 10 Q'd I/Os, 6 read, 4 write
 4W = 4 Q'd I/Os, all write

- Removing QAS simplifies design and validation
 - But ...
 - Reducing impact of **no QAS** requires high queue depths for sequential I/O
 - High queue depths increases I/O latency
 - Heroic seek optimization may not be as effective as expected
 - Will applications be rewritten to make use of large queue depths ?
 - Will early implementations use queue depths > 1 for sequential I/O ?
 - Assuming high queue depths also assumes specific workload profiles - not a general solution
 - This trades 1 problem for a new set

Ultra320 Performance Cost of Changes

Theoretical results - only counts SCSI protocol timings

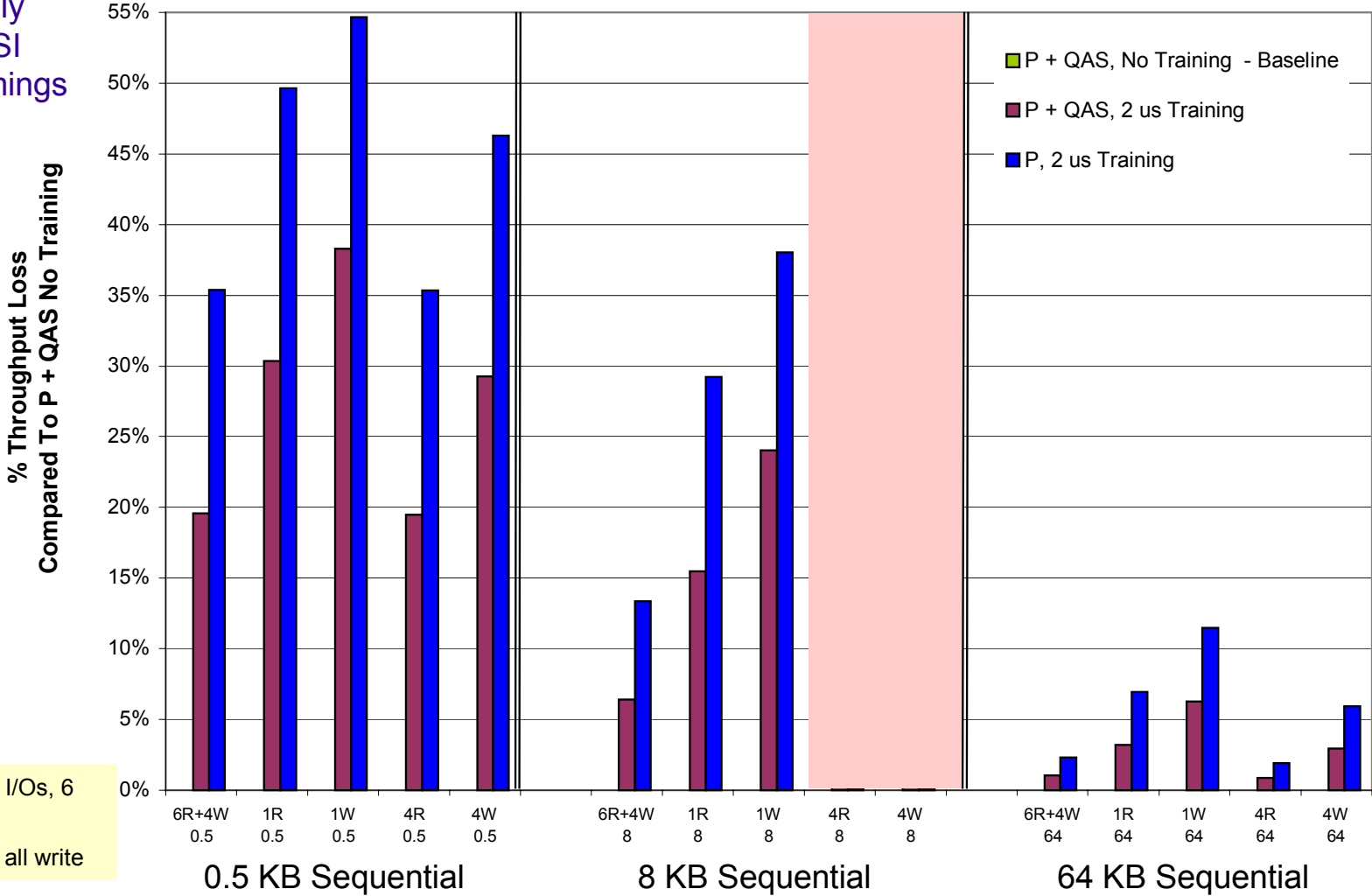


6R+4W = 10 Q'd I/Os, 6 read, 4 write

4W = 4 Q'd I/Os, all write

Ultra640 Performance Cost of Changes

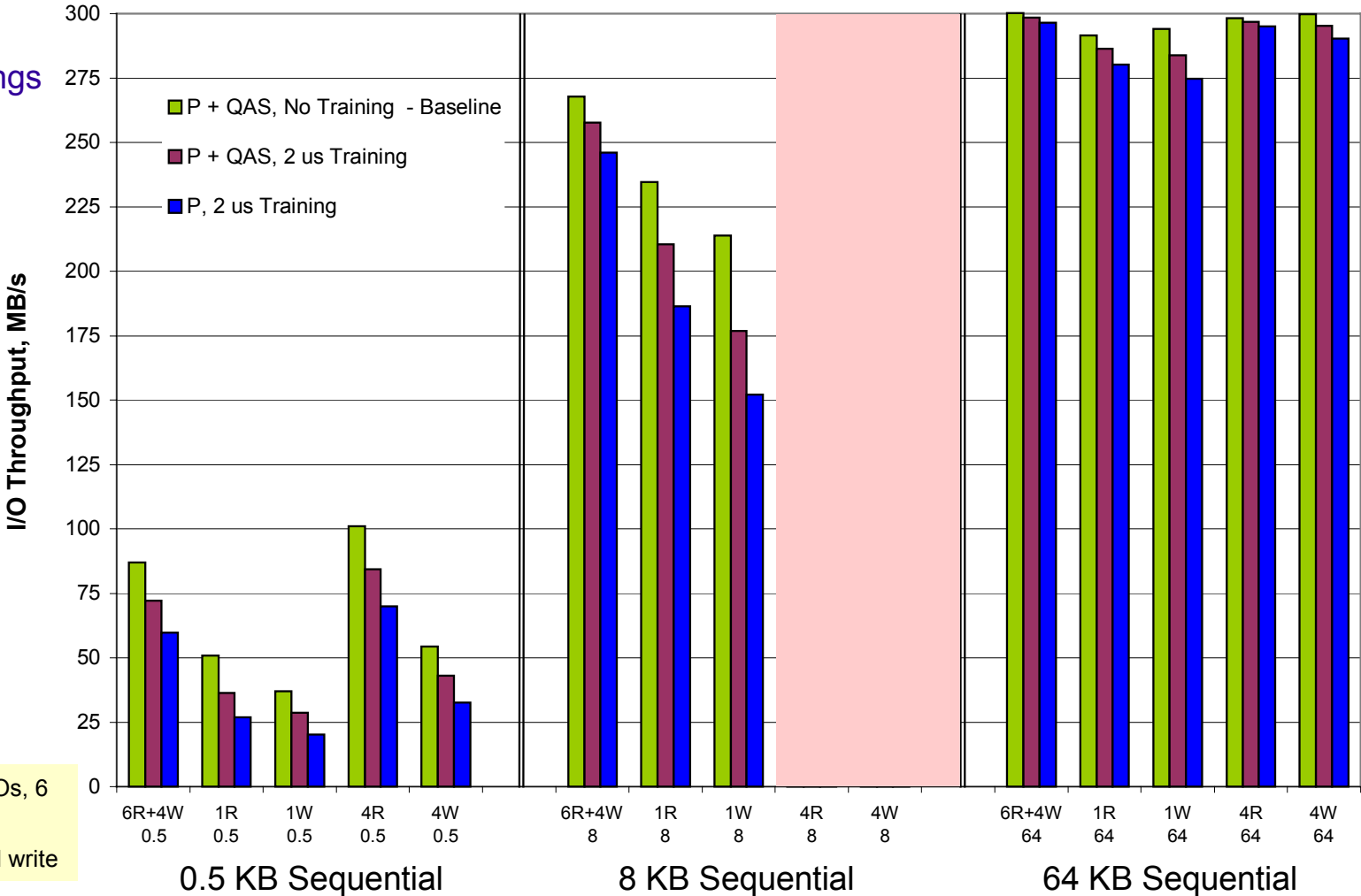
Theoretical results - only counts SCSI protocol timings



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Ultra320 Performance Cost of Changes

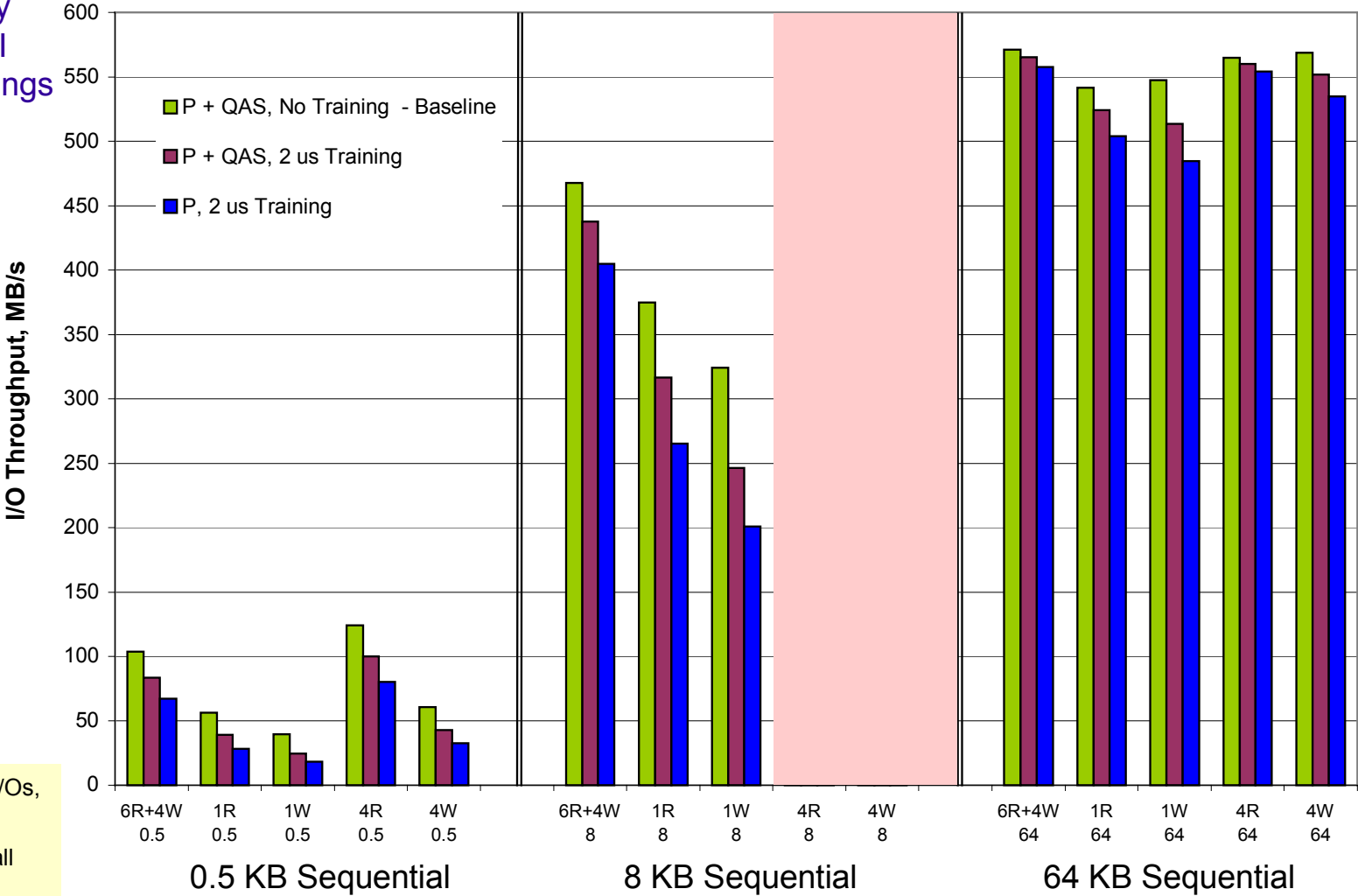
Theoretical results - only counts SCSI protocol timings



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Ultra640 Performance Cost of Changes

Theoretical results - only counts SCSI protocol timings



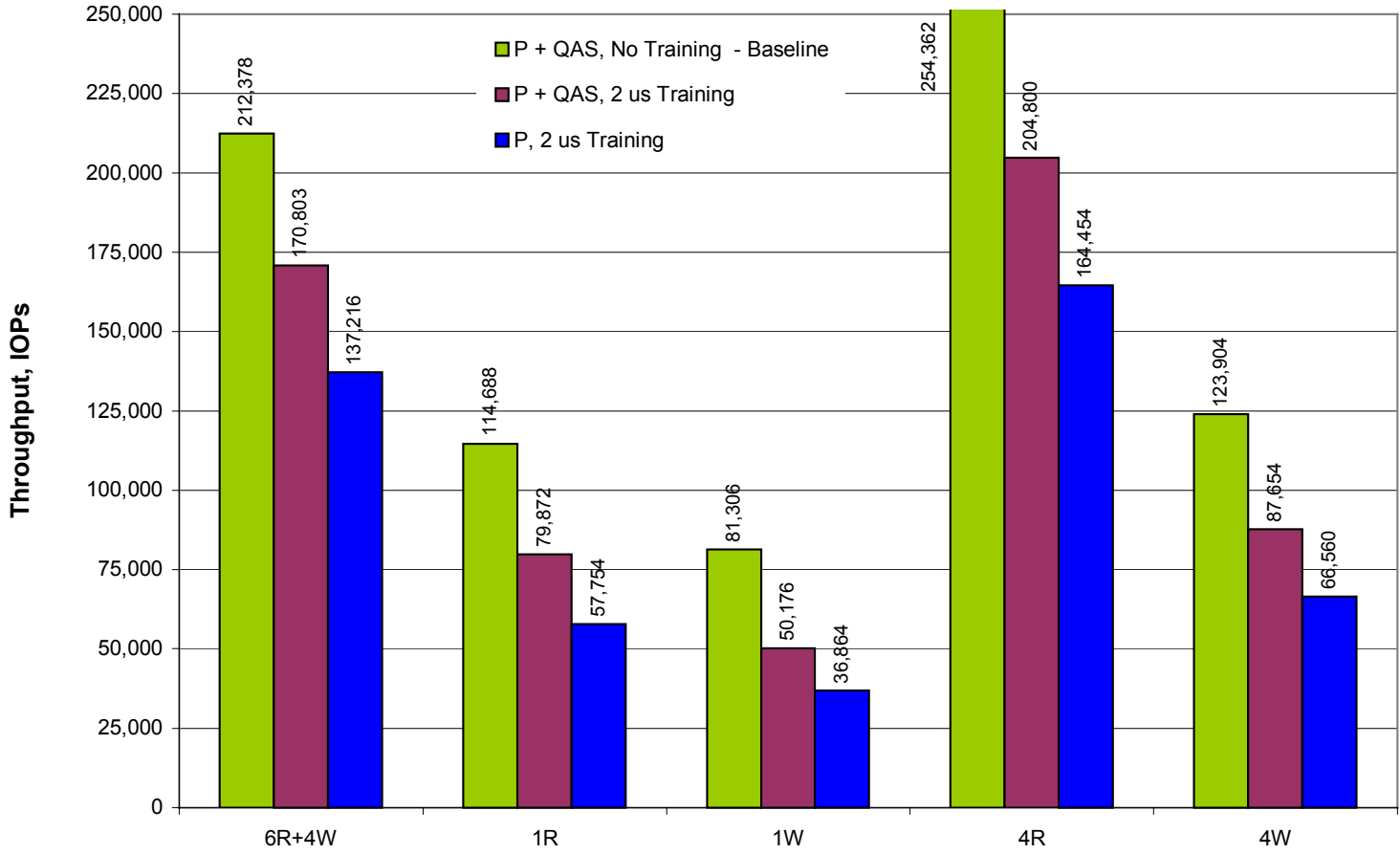
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Ultra640 Performance Cost of Changes, IOPs, 0.5 KB I/Os



6R+4W = 10 Q'd I/Os, 6 read, 4 write
 4W = 4 Q'd I/Os, all write